

# Tutorial 2 : Transitive closure

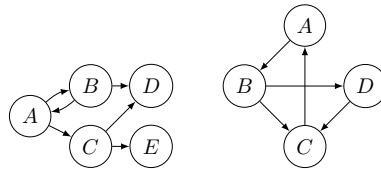
Graph theory, 1st semester.

2022

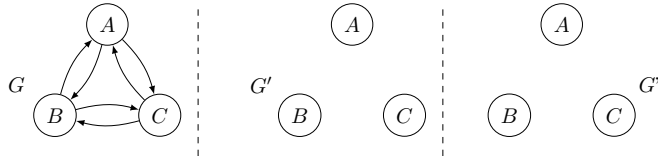
## Exercise 1 — Some examples and properties

Two graphs  $G$  and  $G'$  are  **$\tau$ -equivalent** if  $\tau(G) = \tau(G')$ . A partial graph  $G'$  of  $G$  is  **$\tau$ -minimal**  **$\tau$ -equivalent** to  $G$  if it is  $\tau$ -equivalent to  $G$  and, if we remove an arc of  $G'$ , we get a graph that is not  $\tau$ -equivalent to  $G$ . A graph  $G'$  is  **$\tau$ -minimum  $\tau$ -equivalent** to  $G$  if it is  **$\tau$ -minimal  $\tau$ -equivalent** to  $G$  and minimizes the number of arcs.

1. Draw the transitive closure of the following graphs :



2. Add some arcs to  $G'$  and  $G''$  such that  $G'$  and  $G''$  are  $\tau$ -minimal  $\tau$ -equivalent to  $G$  and  $G''$  contains strictly less arcs than  $G'$ .



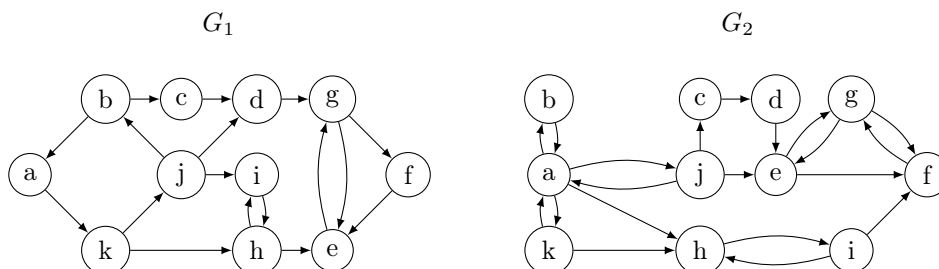
3. Show that  $G$  is strongly connected if and only if the circuit  $(x_1, x_2, \dots, x_n)$  is  $\tau$ -equivalent to  $G$ , where the  $x_i$ s are the nodes of  $G$ .
4. Show that the previous property is false if we replace  $\tau$ -equivalent by  $\tau$ -minimal  $\tau$ -equivalent.
5. Show that an elementary circuit is  $\tau$ -minimum  $\tau$ -equivalent to itself.
6. Show that a graph  $G$  is hamiltonian if and only if, every graph  $\tau$ -minimum  $\tau$ -equivalent to  $G$  is an hamiltonian circuit of  $G$ .

## Exercise 2 — $\tau$ equivalent graph of a acyclic digraph

Let  $G$  be a directed acyclic graph. Show that there exists a unique graph  $G'$  that is  $\tau$ -equivalent to  $G$  and that minimize the number of arcs. Show that  $G'$  is a partial graph of  $G$ .

## Exercise 3 — $\tau$ -equivalent graphs and reduced graphs

Let  $G_1$  and  $G_2$  be the two following graphs :



1. Compute the transitive closures of  $G_1$  and  $G_2$ .
2. Compute the reduced graphs  $G_{1R}$  and  $G_{2R}$  of  $G_1$  and  $G_2$ .
3. Compute the transitive closures of  $G_{1R}$  and  $G_{2R}$ .
4. Show that two graphs are  $\tau$ -equivalent if and only if their reduced graphs are  $\tau$ -equivalent.

**Exercise 4 — Compute the transitive closure of a graph**

1. Describe with pseudo code an algorithm computing the transitive closure of a graph using the classical algorithms enumerating the nodes of a graph. What is its complexity?
2. (a) Let  $A$  be the adjacency matrix of a graph  $G$ . What is  $A^p$  ?  
 (b) Let  $B$  be the adjacency matrix of a graph  $G$  with 1 on the diagonal cells. What is  $B^p$  ?  
 (c) Describe in pseudocode an algorithm computing the transitive closure of a graph using the powers of either the matrix  $A$  or the matrix  $B$ . What is its complexity?
3. The Roy Warshall algorithm is the following :

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**Require:** A directed graph  $G = (U, A)$

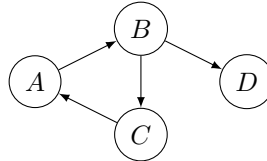
**Ensure:** The transitive closure of  $G$

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for  $w \in U$  do
  for  $u \in U$  do
    for  $v \in U$  do
      if  $(u, w) \in A$  and  $(w, v) \in A$  then Add  $(u, v)$  to  $A$ .
  
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- (a) Run the algorithm on the following graph :



- (b) What is its complexity ?  
 (c) Describe the algorithm using only the adjacency matrix of  $G$ .  
 (d) Run the algorithm on the following matrix.

$$\begin{pmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{pmatrix}$$

**Exercise 5 — Equipment of a workshop**

Nine machines,  $a, b \dots, h$  and  $i$ , are installed in a workshop. Mechanical parts are produced by successive machines (drills, polishers, welds, ...). We want to build conveyor belts at minimum cost in order to move the pieces from a machine to another. The following array gives, for each machine  $M$ , which machines may be used in a next step to build a product. We must link  $M$  to each of those machines by a succession of belts. Model this problem with a graph problem and solve it.

$a$	$b, c, d, e, f, g, h, i$
$b$	$a, c, d, e, f, g, h, i$
$c$	$d, e$
$d$	$e$
$e$	$d$
$f$	$d, e, g, h, i$
$g$	$d, e, f, h, i$
$h$	$d, e, f, g, i$
$i$	$d, e, f, g, h$